

Systematic API Testing through Model Checking and Executable Contracts

TLA+ Community Meeting

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NOVA University Lisbon

April 12, 2026



Cracking the Ice:

How to test an iceberg?

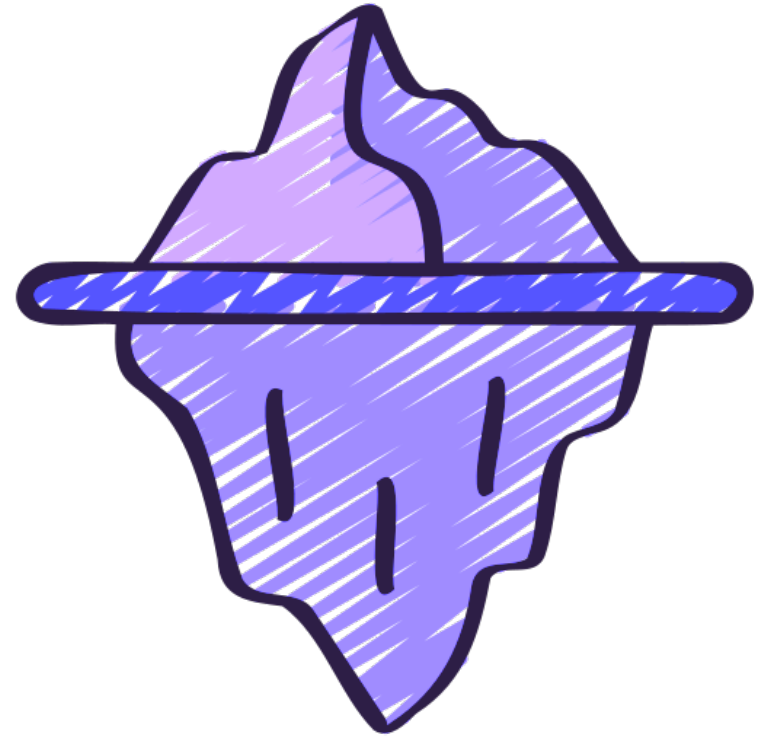


The Iceberg *aka "The SUT"*

Implementation

Heterogeneous

Black-box



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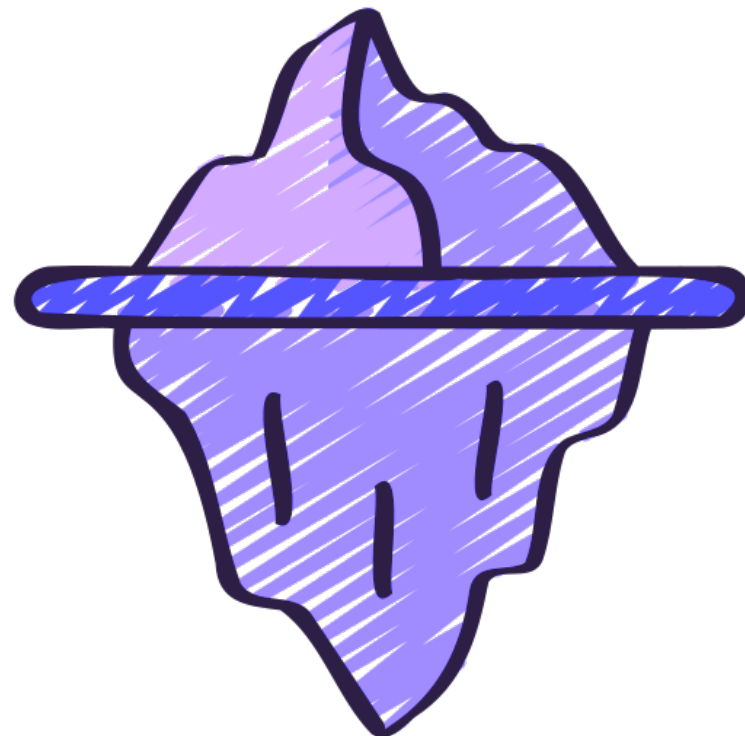
API

RESTful

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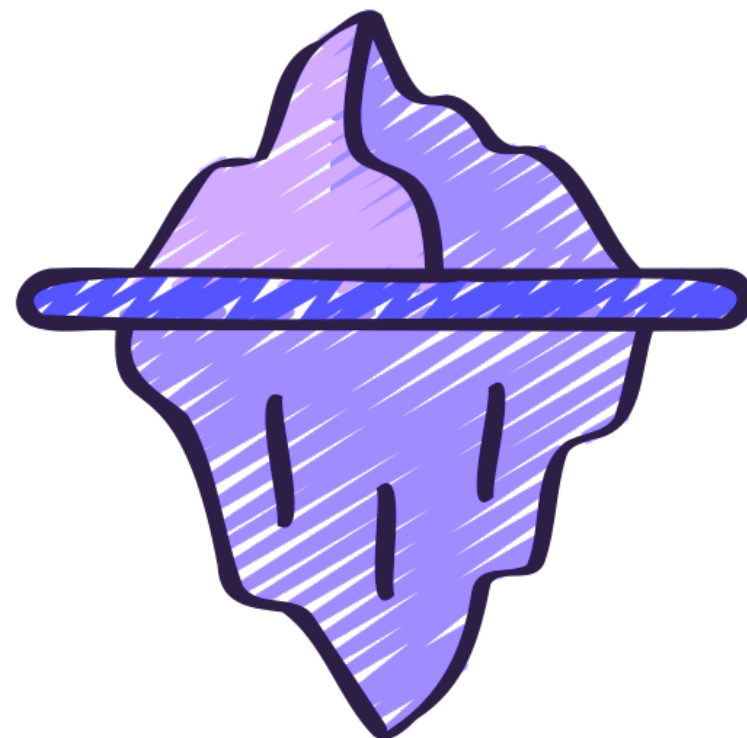
The Iceberg *aka "The SUT"*

API

RESTful

REpresentational State Transfer (REST)

Guidelines on building web services on top of HTTP.

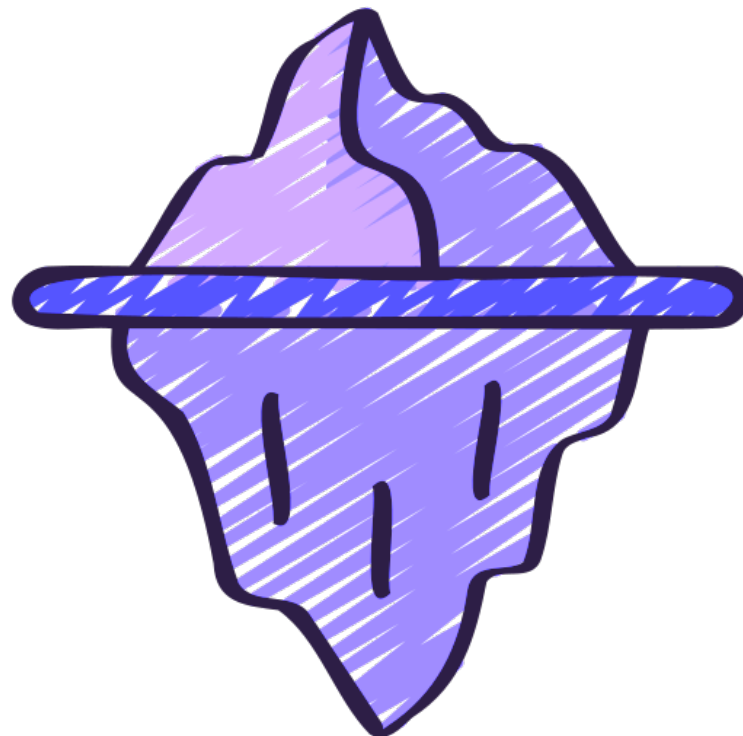


The Iceberg *aka "The SUT"*

API

RESTful

Described in
OAS



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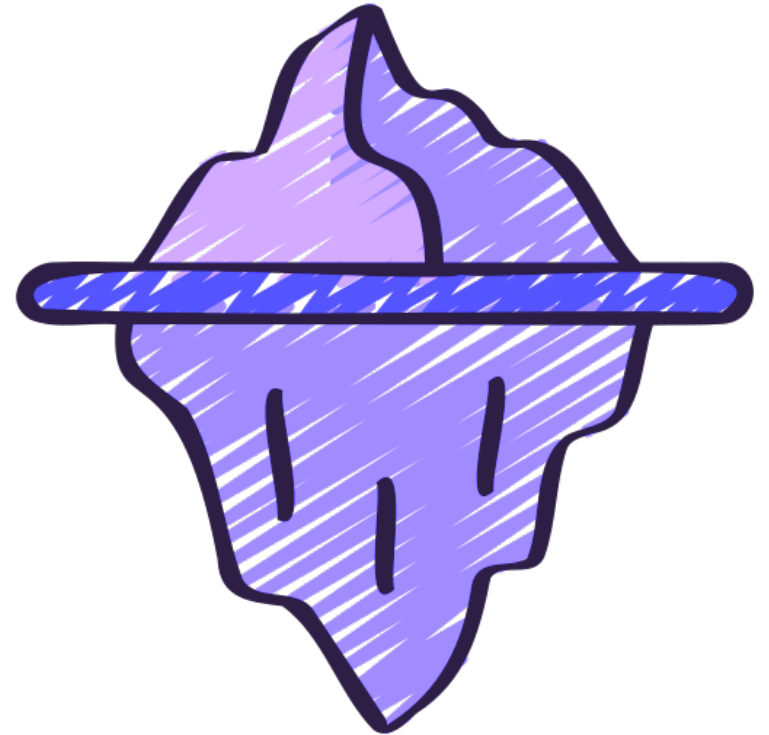
API

RESTful

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OAS (Open API Specification)

Standard, language-agnostic format for specifying HTTP APIs.



The Iceberg *aka "The SUT"*

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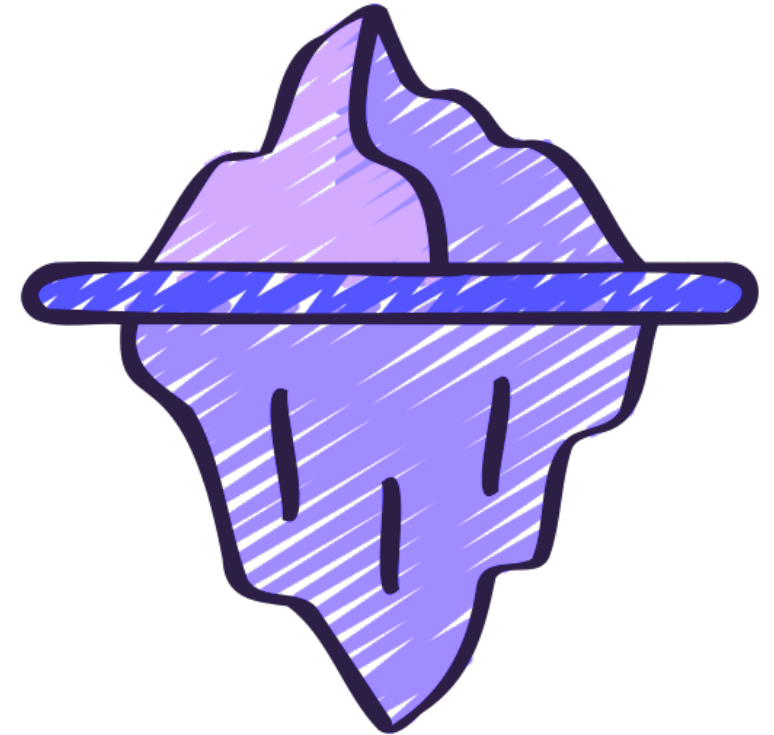
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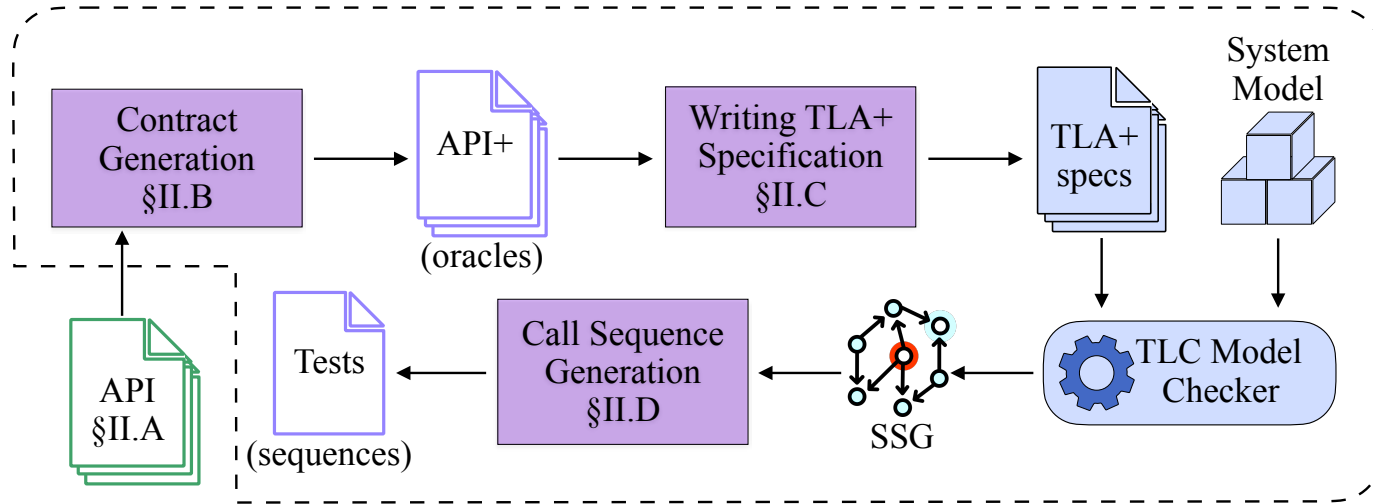
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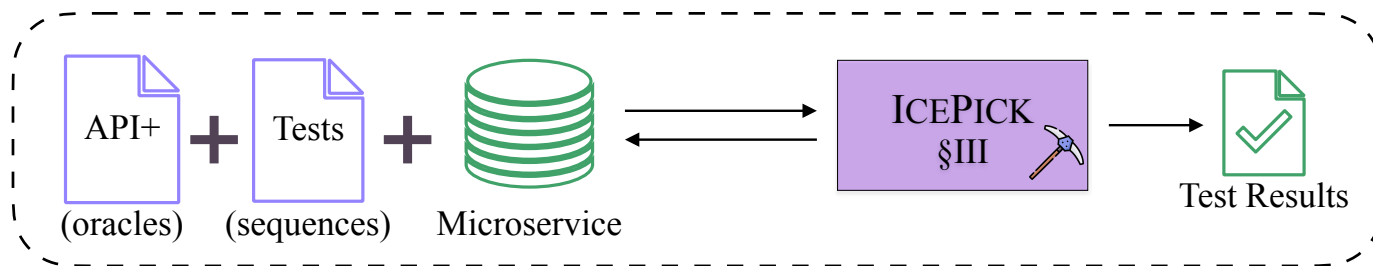


Solution – Overview

ICEPICK Phase 1 - Specification pre-processing

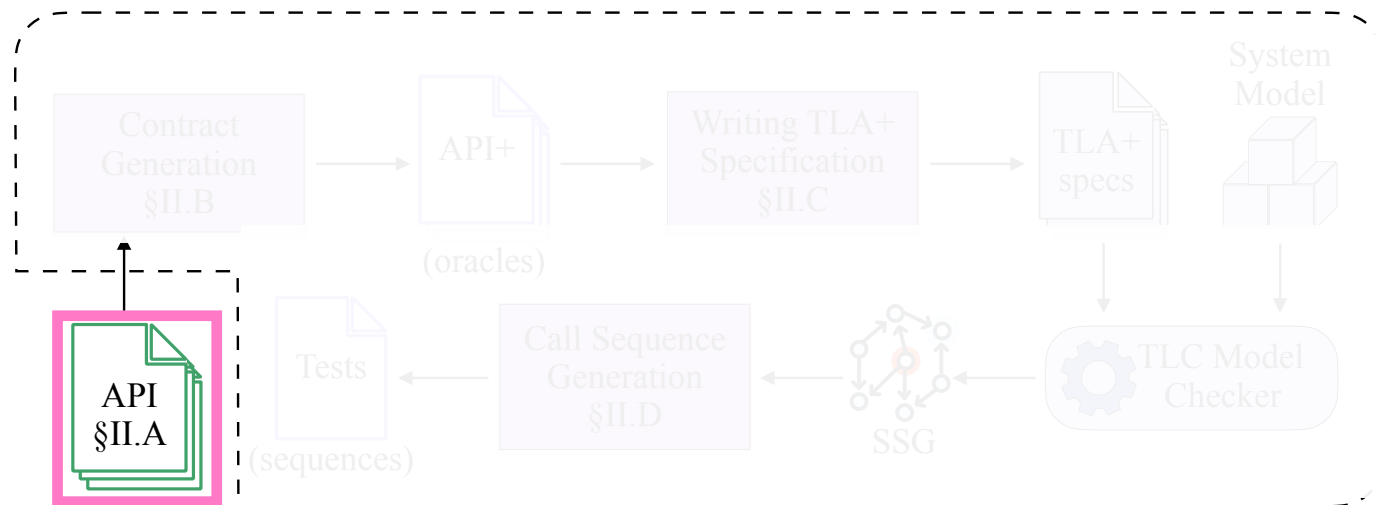


ICEPICK Phase 2 - Testing

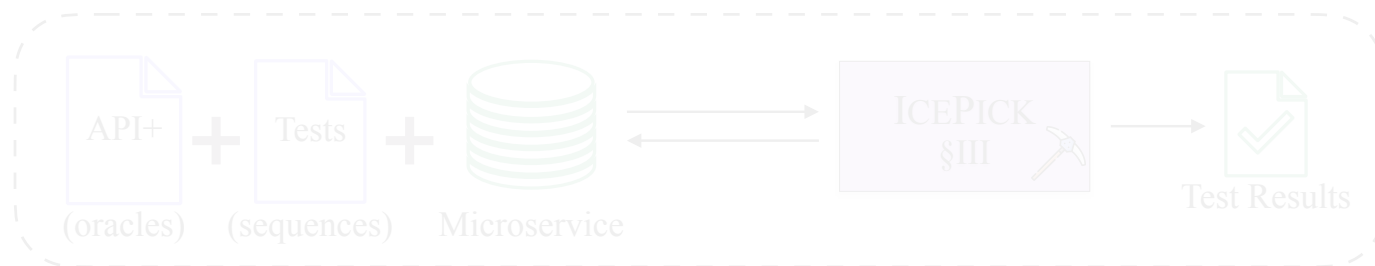


Phase 1 – The Standard

ICEPICK Phase 1 - Specification pre-processing



ICEPICK Phase 2 - Testing



OpenAPI Specification

OAS offers:

- operation endpoints
- expected data types
- scenario-based response codes

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Running Example

Tournament

- tid
- capacity (c)
- players (ps)

Player

- pid
- name
- tournaments (ts)

Enrolments

- eid
- pid
- tid

OpenAPI Specification

OAS offers:

- operation endpoints
- expected data types
- scenario-based response codes

```
/players:  
  post:  
    summary: Add a new player.  
    operationID: postPlayer  
    requestBody:  
      content:  
        application/json:  
          schema:  
            $ref: #/schemas/Player  
            required: true  
    responses:  
      200:  
        description: Successful insertion.  
        content:  
          application/json:  
            schema:  
              $ref: #/schemas/Player
```

OpenAPI Specification – Not enough?

OAS offers:

- operation endpoints
- expected data types
- scenario-based response codes

Existing solutions derive tests exclusively from OAS specifications.

- poorly written (param. naming)
- incomplete (op. dependencies)
- outdated (evolution)

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OAS offers:

- operation endpoints
- expected data types
- scenario-based response codes

Existing solutions derive tests exclusively from OAS specifications.

- poorly written (param. naming)
- incomplete (op. dependencies)
- outdated (evolution)
- status-code oracles (**5xx ISE?**)

```
responses:  
  500:  
    description: Unexpected error.  
    content:  
      application/json:  
        schema:  
          $ref: "#/schemas/ErrorObject"
```

OpenAPI Specification – Not enough?

OAS offers:

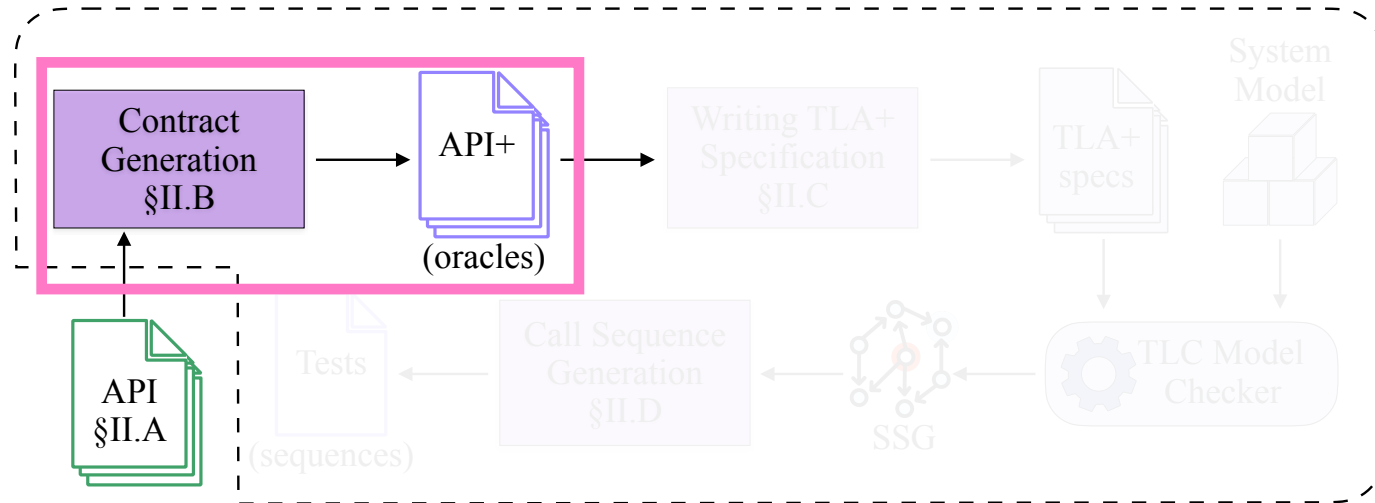
- operation endpoints
- expected data types
- scenario-based response codes

OAS lacks detail to capture system behaviour

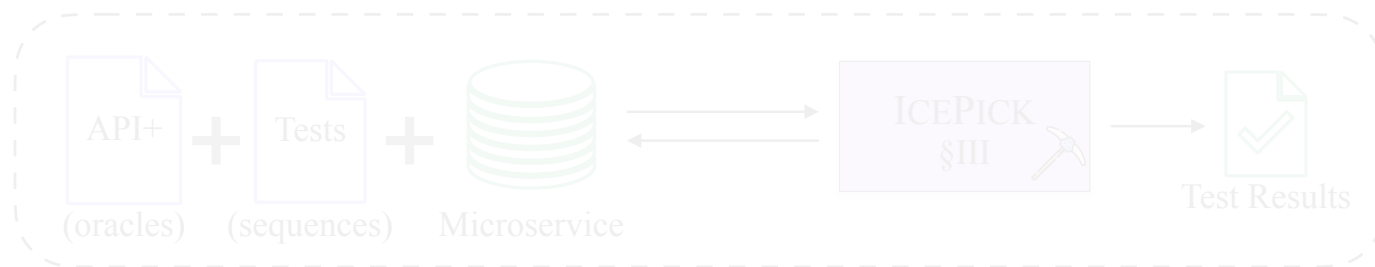
- preconditions
- postconditions
- invariants

Phase 1 – Extending the Standard

ICEPICK Phase 1 - Specification pre-processing

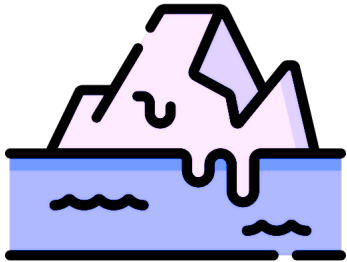


ICEPICK Phase 2 - Testing



GLACIER – Yet another specification language?

First-order logic based
contract specification
language for RESTful APIs

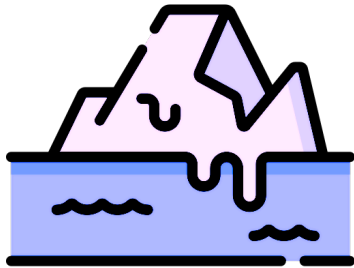


Enriches API specifications
with practical properties for
test generation

- Aims to solve the oracle problem, by adding
 - **preconditions**
 - **postconditions**
 - **invariants**
- RESTful API tailored constructs:
 - `res_code, res_body / req_body`
- Useful spec constructs
 - `@, prev`

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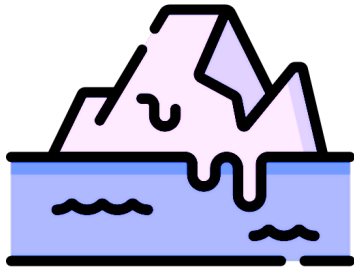


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```
/players:  
  POST:  
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    requires:  
      - res_code(GET /players/req_body(@){pid}) = 404  
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      - req_body(@) = res_body(@)
```

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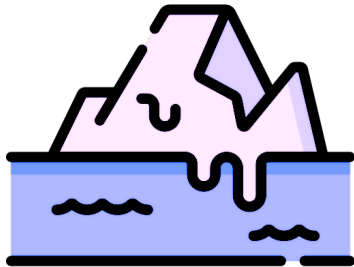


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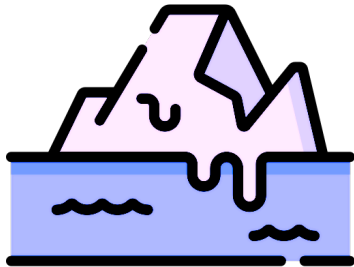


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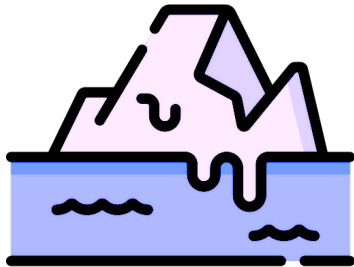


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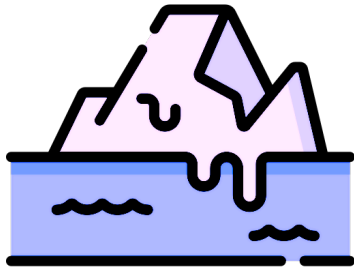


Enriches API specifications
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```
/players/{pid}:  
DELETE:  
  operationID: deletePlayer  
  requires:  
  - res_code(GET /players/{pid}) = 200  
  ensures:  
  - res_code(GET /players/{pid}) = 404  
  - req_body(@) = prev(res_body(GET /players/{pid}))
```

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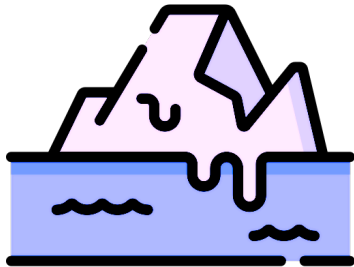


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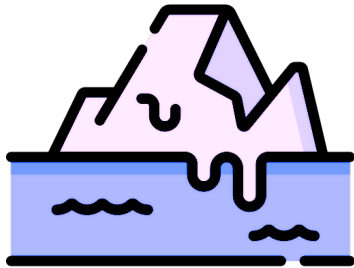


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Invariants?

GLACIER – Yet another specification language?

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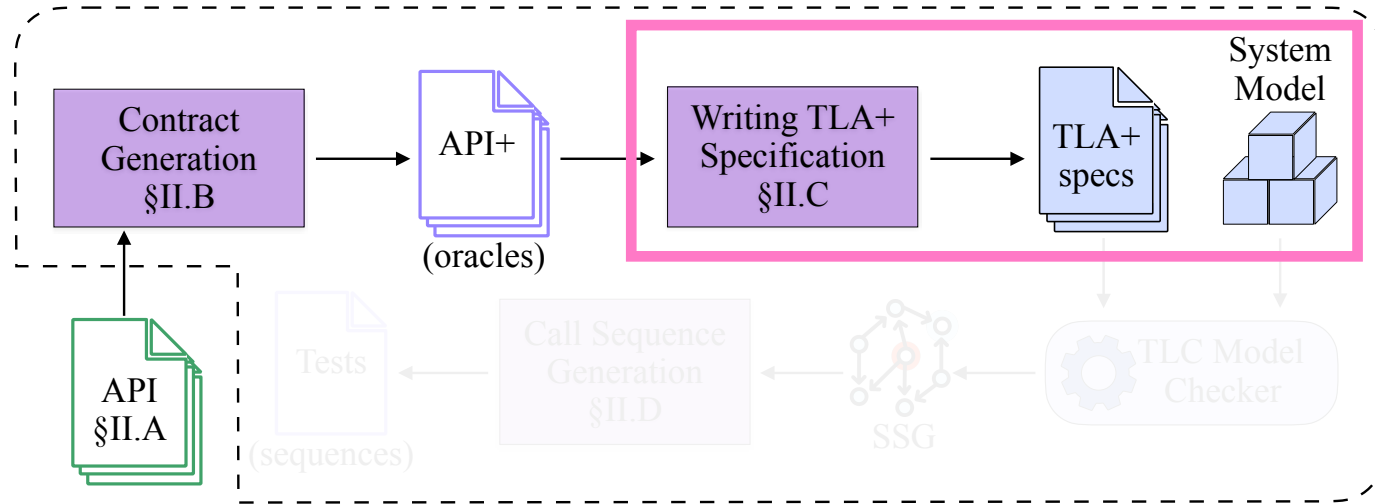
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`invariants:`

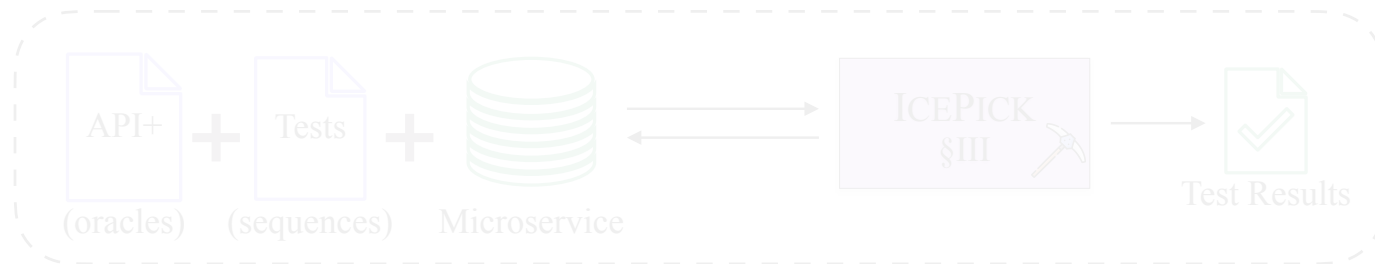
```
- for t in resp_body(GET /tournaments) :-  
  res_body(GET /tournaments/{t.tid}/players).len ≤  
  res_body(GET /tournaments/{t.tid}/capacity)
```

Phase 1 – The TLA+ Specification

ICEPICK Phase 1 - Specification pre-processing



ICEPICK Phase 2 - Testing



Phase 1: TLA+ Specification – The Abstraction

- Life-cycle oriented
 - create → update → remove

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- Life-cycle oriented
 - create → update → remove
- Focus on operations inducing state transitions
 - post
 - delete
- Read-only operations are not considered
 - get
- Update operations modifying abstracted fields are handled at a later stage
 - put
 - patch

Phase 1: Tournaments API – TLA+ Specification

CONSTANTS PID, TID, EID, N

Resource identifiers and the Naturals set.

Phase 1: Tournaments API – TLA+ Specification

```
CONSTANTS PID, TID, EID, N
```

Resource identifiers and the Naturals set.

```
ASSUME  
  ^ PID ≠ {}  
  ^ TID ≠ {}  
  ^ EID ≠ {}  
  ^ N ∈ SUBSET (Nat \ {0})  
  ^ IsFiniteSet(N)
```

All identifier sets are non-empty and N is a finite set of positive natural numbers.

Phase 1: Tournaments API – TLA+ Specification

```
PlayerType      = [ts: SUBSET TID]
TournamentType = [ps: SUBSET PID, c: N]
EnrolmentType  = [pid: PID, tid: TID]
```

Resources are represented as records, omitting schema fields irrelevant to reachability.

Phase 1: Tournaments API – TLA+ Specification

```
PlayerType      = [ts: SUBSET TID]  
TournamentType = [ps: SUBSET PID, c: N]  
EnrolmentType  = [pid: PID, tid: TID]
```

Resources are represented as records, omitting schema fields irrelevant to reachability.

```
VARIABLES players, tournaments, enrolments, final
```

One variable per resource type and a terminal state flag.

Phase 1: Tournaments API – TLA+ Specification

```
TypeInv =  
  ∧ IsMap(PID, PlayerType, players)  
  ∧ IsMap(TID, TournamentType, tournaments)  
  ∧ IsMap(EID, EnrolmentType, enrolments)  
  ∧ final ∈ BOOLEAN
```

Resource maps are keyed by their identifiers and `final` is a boolean.

Phase 1: Tournaments API – TLA+ Specification

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TypeInv =  
  ∧ IsMap(PID, PlayerType, players)  
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  ∧ final ∈ BOOLEAN
```

Resource maps are keyed by their identifiers and `final` is a boolean.

```
Init =  
  ∧ players = EmptyMap  
  ∧ tournaments = EmptyMap  
  ∧ enrolments = EmptyMap  
  ∧ final = FALSE
```

Initial state sets all resource maps to empty and `final` to `FALSE`.

Phase 1: Tournaments API – TLA+ Specification

```
Inv1 =  
  ∀ tid ∈ TID : tid ∈ DOMAIN tournaments ⇒  
    Cardinality(tournaments[tid].ps) ≤ tournaments[tid].c
```

GLACIER invariants become TLC-checked state invariants.

Phase 1: Tournaments API – TLA+ Specification

```
Requires(pid) = pid ∉ DOMAIN players  
Ensures(pid)  = pid ∈ DOMAIN players
```

Operational contracts become action guards and next-state constraints.

Phase 1: Tournaments API – TLA+ Specification

```
postPlayer(pid) =  
  ∧ Requires(pid)  
  ∧ players' = MapPut(players, pid, [ts ↦ {}])  
  ∧ Ensures(pid)  
  ∧ final' = FinalState(players', tournaments, enrolments)  
  ∧ UNCHANGED<<tournaments, enrolments>>
```

Each state-mutating API operation → TLA+ action.

Phase 1: Tournaments API – TLA+ Specification

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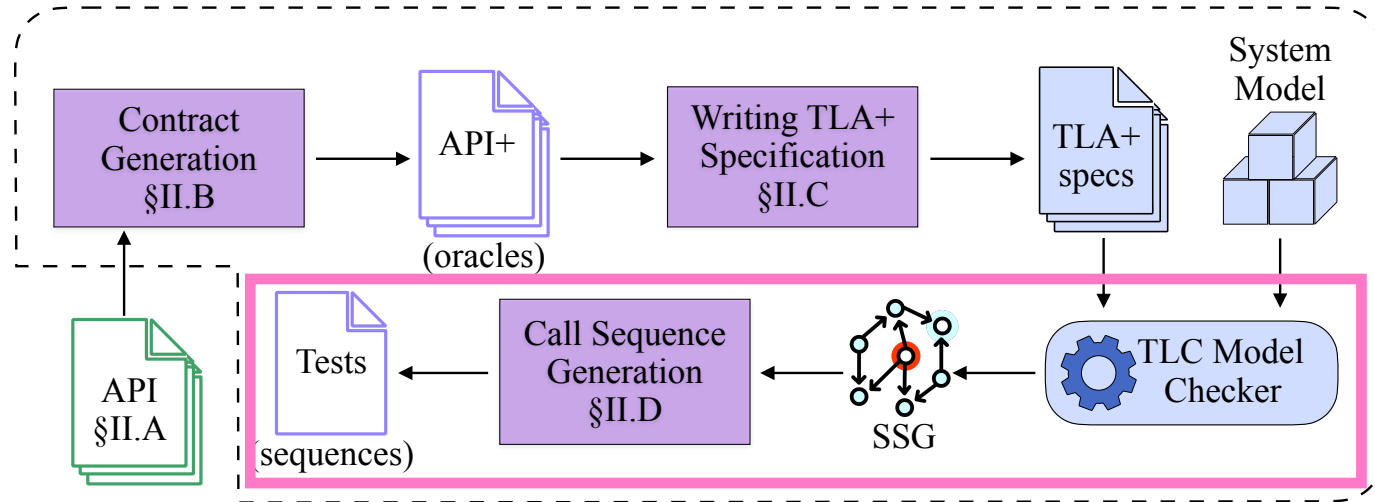
Each state-mutating API operation → TLA+ action.

```
FinalState(players, tournaments, enrolments) =  
  ∀ pid ∈ PID, tid ∈ TID, eid ∈ EID :  
    ∧ pid ∈ DOMAIN players  
    ∧ tid ∈ DOMAIN tournaments  
    ∧ eid ∈ DOMAIN enrolments
```

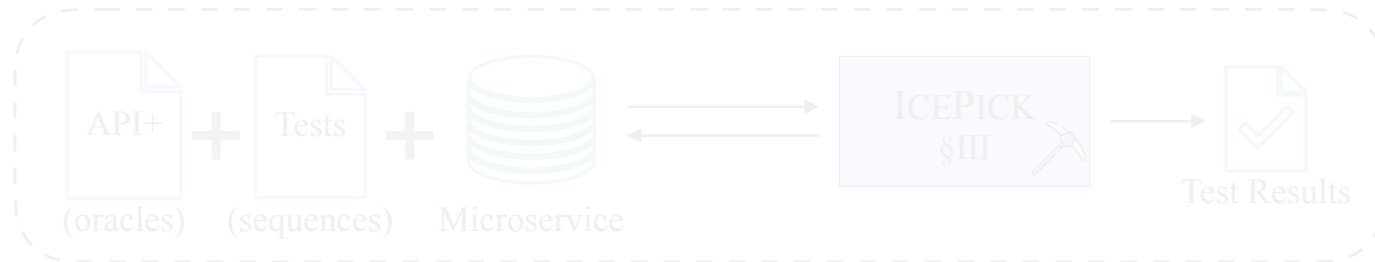
We reach a final state when all resources (model values) have been explored.

Phase 1 – Call Sequence Generation

ICEPICK Phase 1 - Specification pre-processing



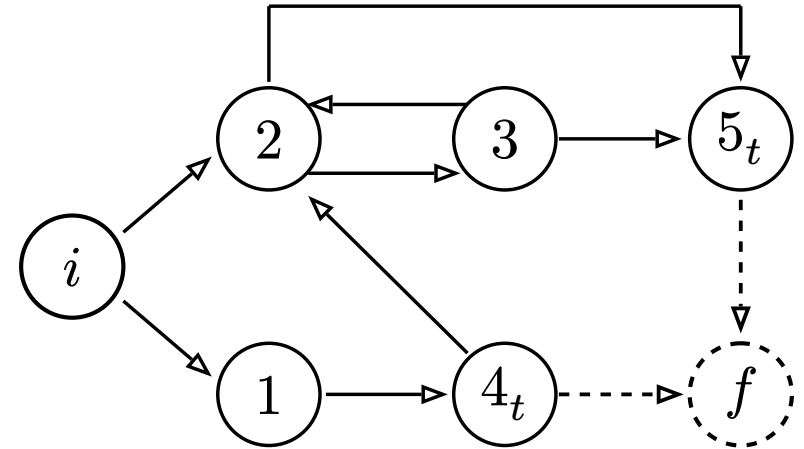
ICEPICK Phase 2 - Testing



Phase 1 – The State Space Graph (SSG)

The SSG is a directed graph $G = (V, E)$, where:

- V is the set of states explored by TLC
 - single initial state i
 - multiple terminal states
 - final sink state f
 - For every terminal state t , $(t, f) \in E$
- E is the set of state transitions (API operations)
 - The graph has no parallel edges



Phase 1 – Call Sequence Generation

Custom **coverage-guided** algorithm based on breadth-first search (BFS).

The algorithm works in three stages:

1. Collect paths from i , classifying them as complete (reaching f) or incomplete
2. Collect a path from each state to f
3. Complete the incomplete paths

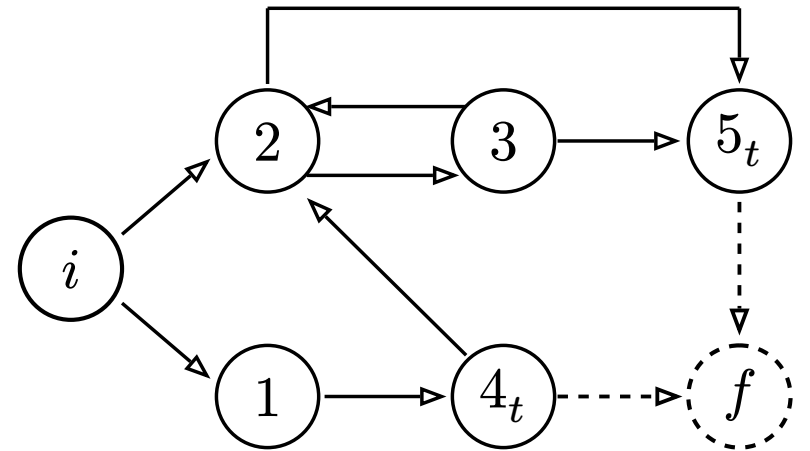
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Phase 1 – Call Sequence Generation (Stage 1)

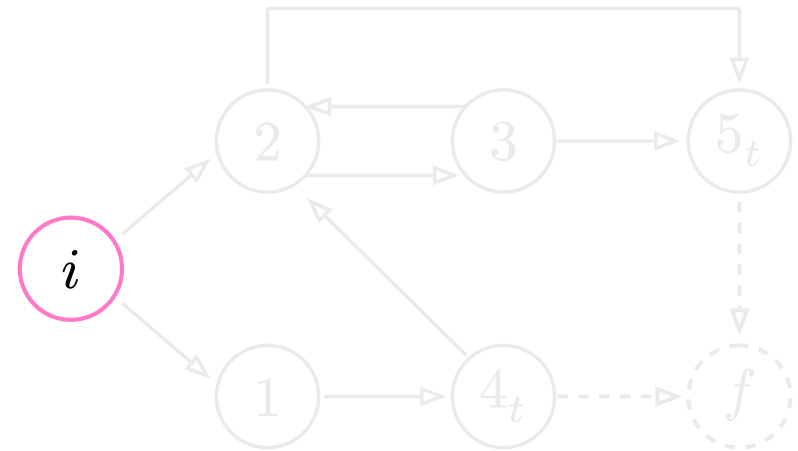


Phase 1 – Call Sequence Generation (Stage 1)

i $\langle i \rangle$

found:

T						T
i	1	2	3	4_t	5_t	f

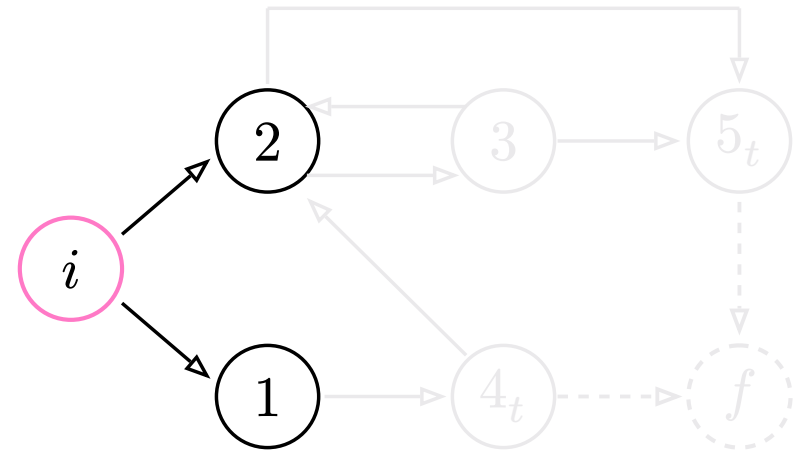
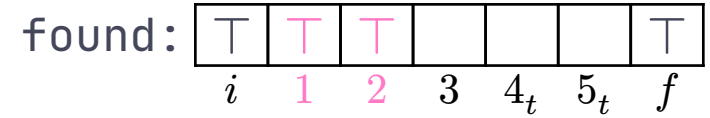
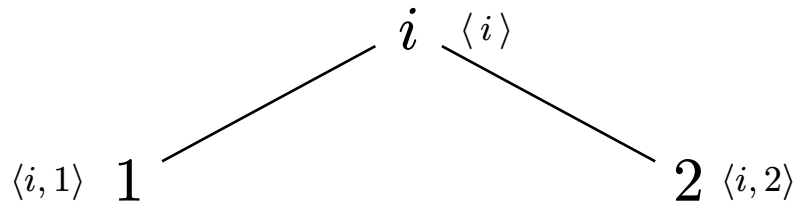


fifo: [i]

cmp : {}

inc : {}

Phase 1 – Call Sequence Generation (Stage 1)



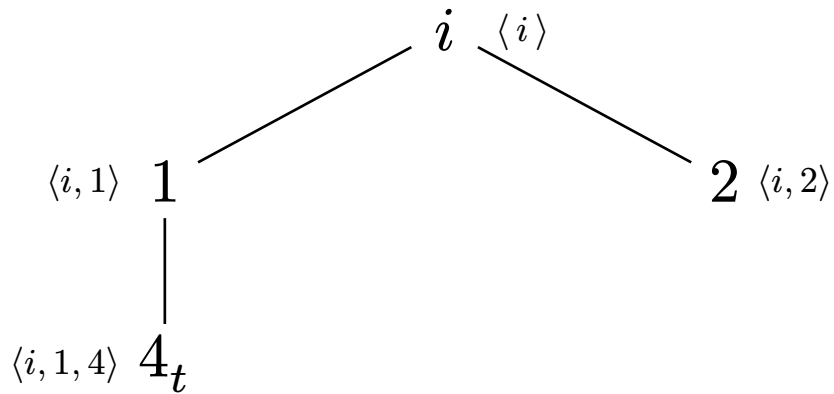
fifo: [*i*, 1, 2]
cmp : {}
inc : {}

Expanding state *i*

We found two new states: 1 and 2

Mark as found; Add to fifo

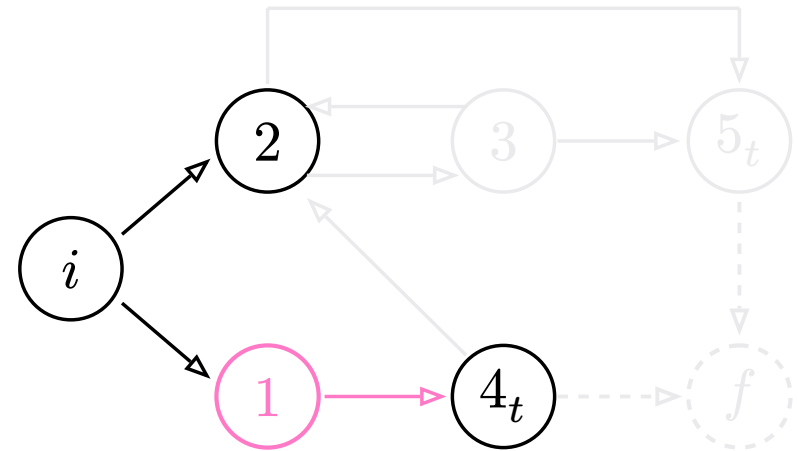
Phase 1 – Call Sequence Generation (Stage 1)



fifo: [*i*, 1, 2, 4_{*t*}]
cmp : { }
inc : { }

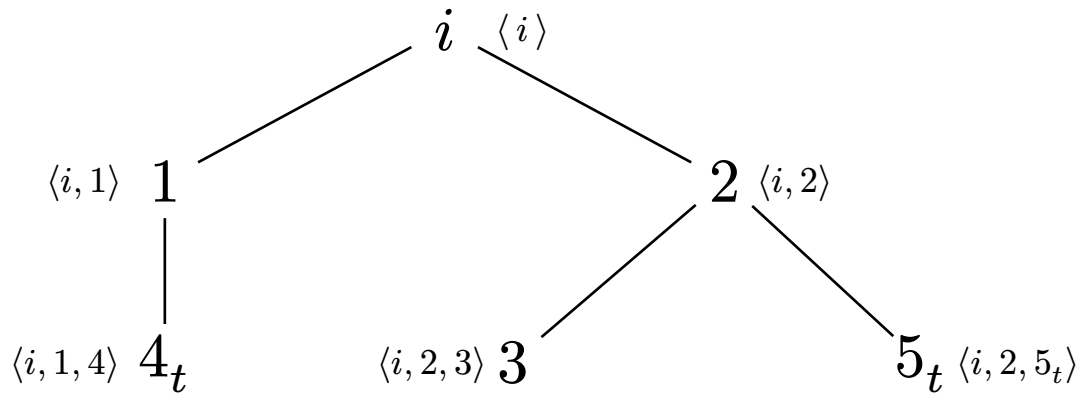
found:

T	T	T		T		T
<i>i</i>	1	2	3	4 _{<i>t</i>}	5 _{<i>t</i>}	<i>f</i>



Expanding state 1
Found a new state: 4_{*t*}
Mark as found. Add to fifo

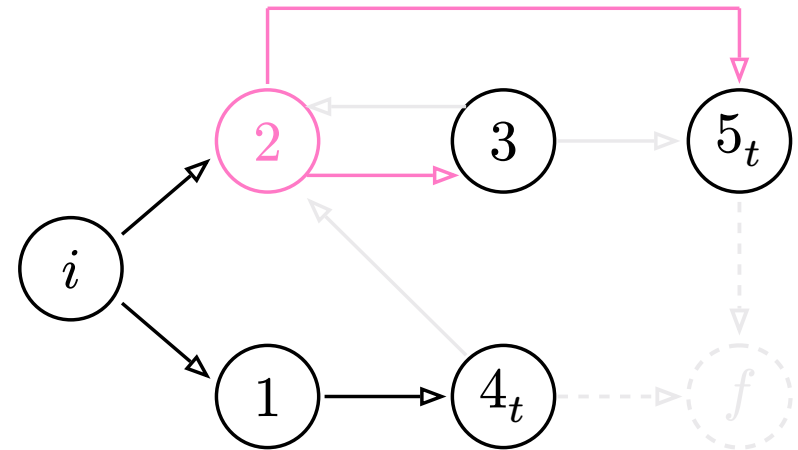
Phase 1 – Call Sequence Generation (Stage 1)



fifo: [i , 1, 2, 4_t , 3, 5_t]
cmp : { }
inc : { }

found:

T	T	T	T	T	T	T
i	1	2	3	4_t	5_t	f

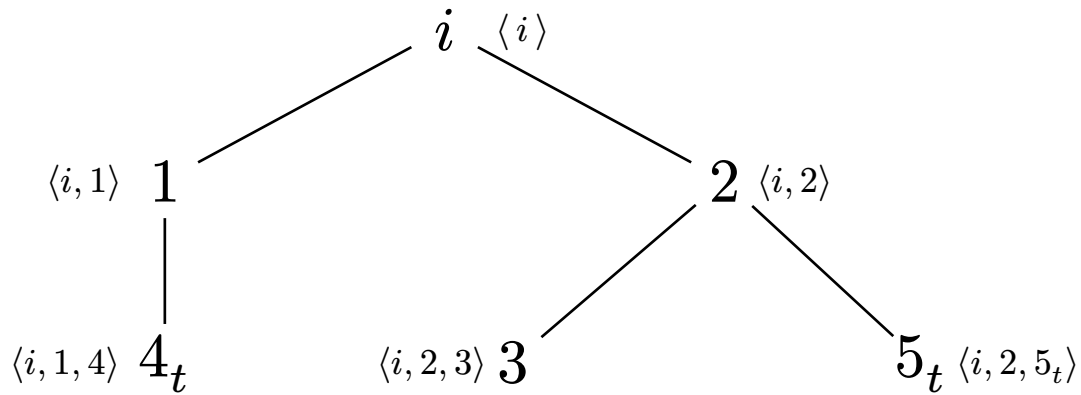


Expanding state 2

Found two new states: 3 and 5_t

Mark as found. Add to fifo

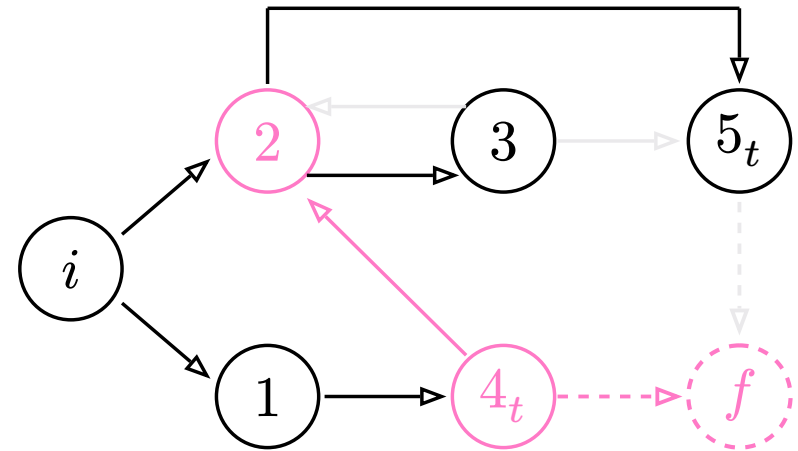
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fifo: [i , 1, 2, 4_t , 3, 5_t]
cmp : {}
inc : {}

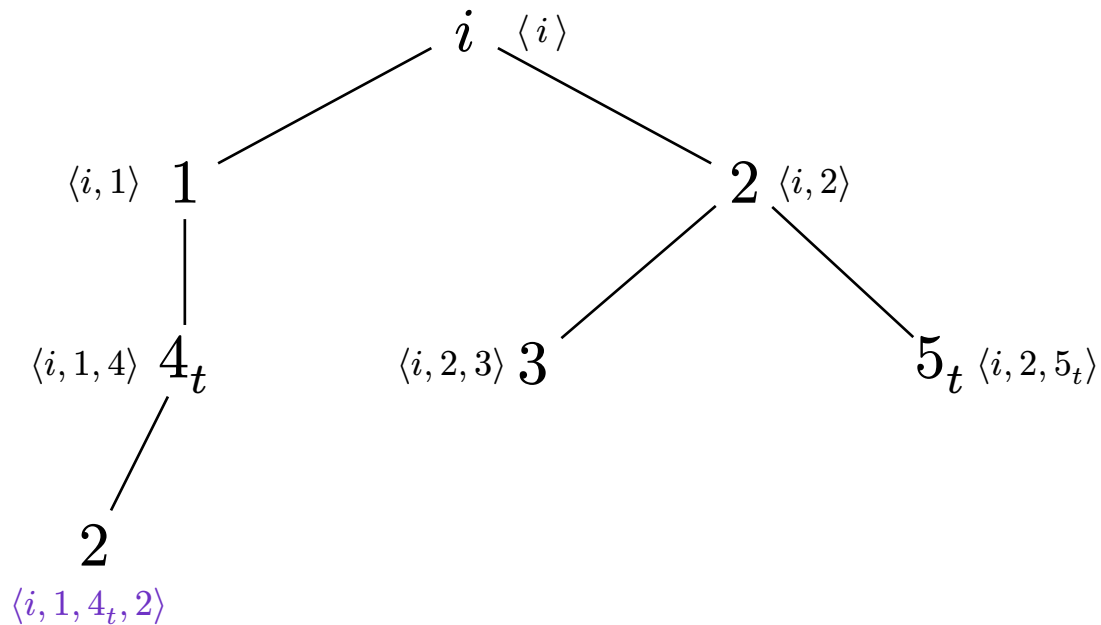
found:

T	T	T	T	T	T	T
i	1	2	3	4_t	5_t	f

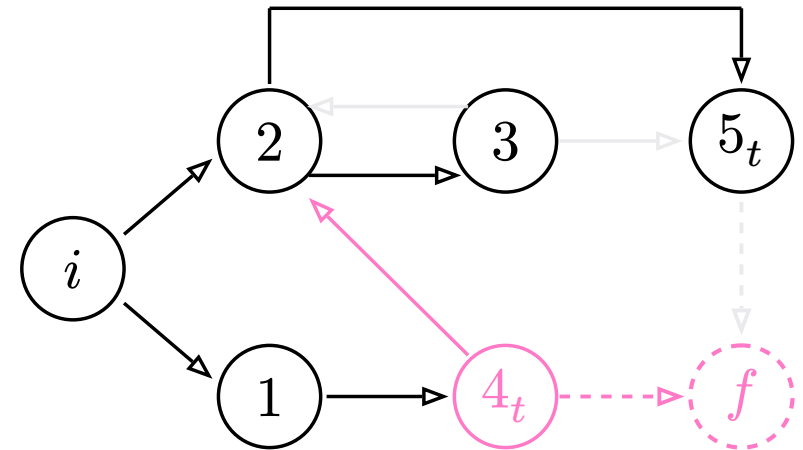
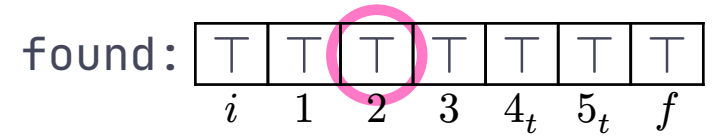


Expanding state 4_t

Phase 1 – Call Sequence Generation (Stage 1)

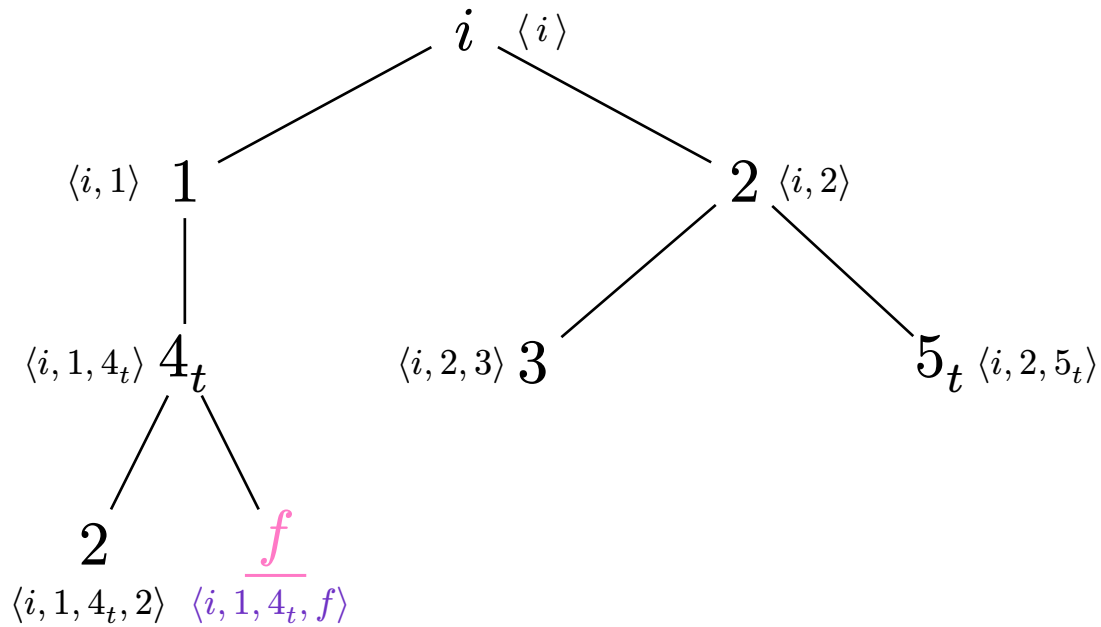


fifo: [i , 1, 2, 4_t , 3, 5_t]
cmp : { }
inc : { $\langle i, 1, 4_t, 2 \rangle$ }

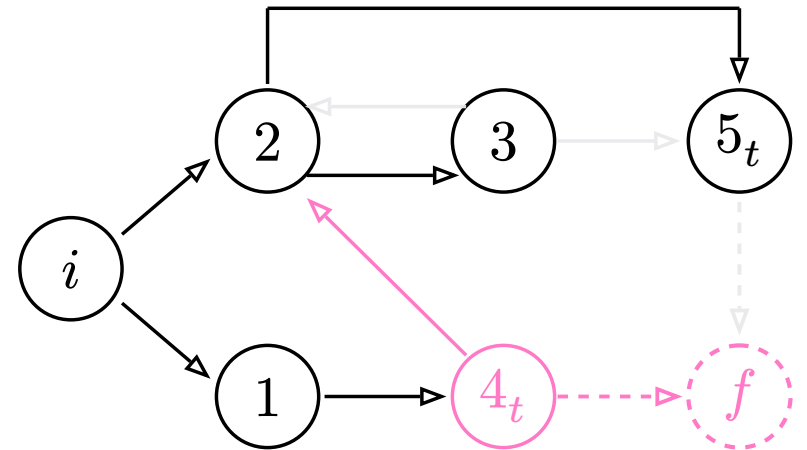
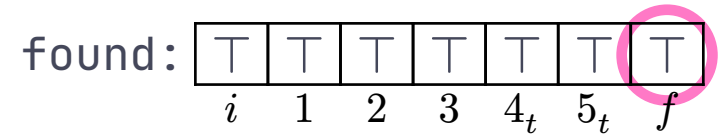


Expanding state 4_t
State 2 has been found
Add path to **inc**

Phase 1 – Call Sequence Generation (Stage 1)

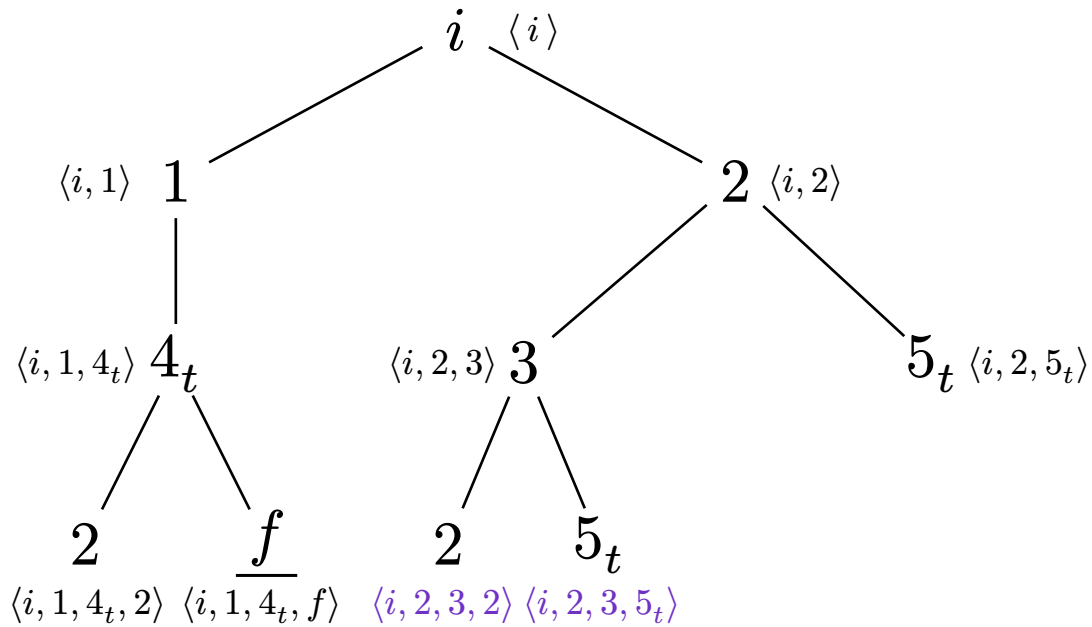


fifo: $[i, 1, 2, 4_t, 3, 5_t]$
 cmp : $\{ \langle i, 1, 4_t, f \rangle \}$
 inc : $\{ \langle i, 1, 4_t, 2 \rangle \}$



Expanding state 4_t
 Final state f has been found
 Add path to **cmp**

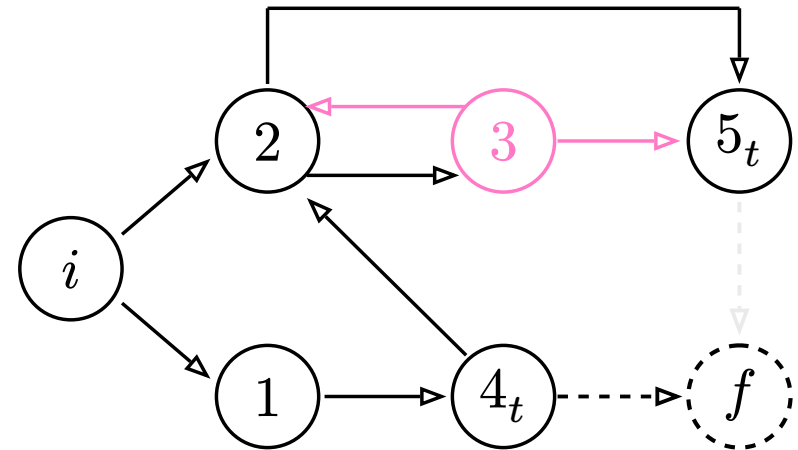
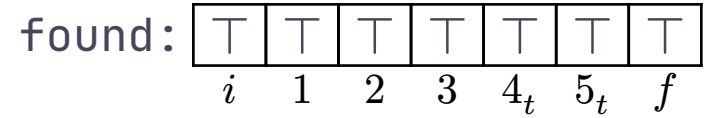
Phase 1 – Call Sequence Generation (Stage 1)



fifo: [i , 1, 2, 4_t , 3, 5_t]

cmp : { $\langle i, 1, 4_t, f \rangle$ }

inc : { $\langle i, 1, 4_t, 2 \rangle$, $\langle i, 2, 3, 2 \rangle$, $\langle i, 2, 3, 5_t \rangle$ }

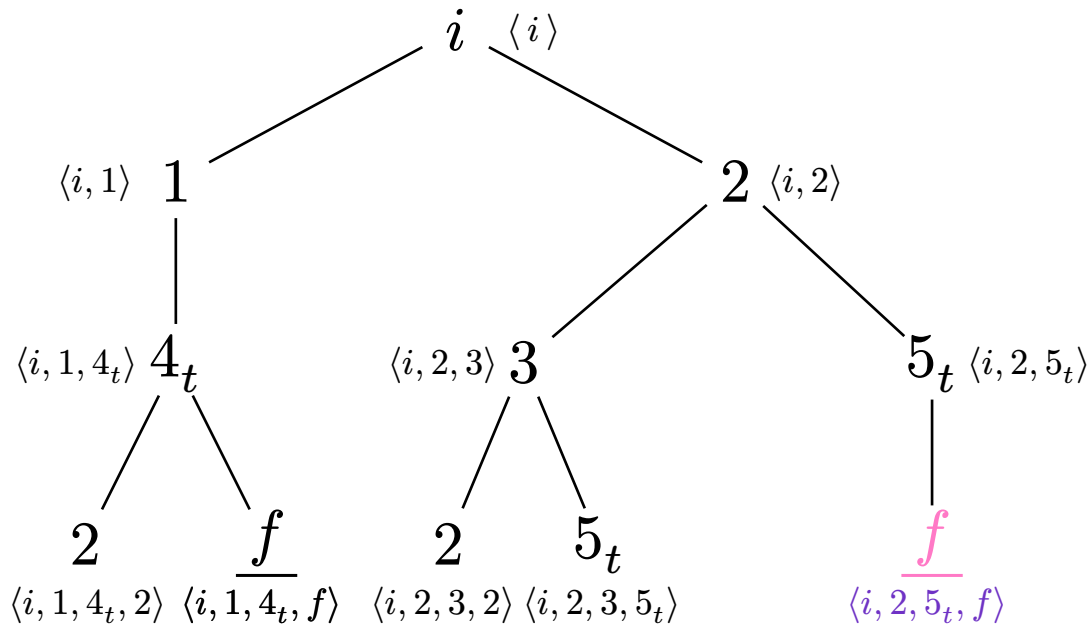


Expanding state 3

No new states

Add paths to **inc**

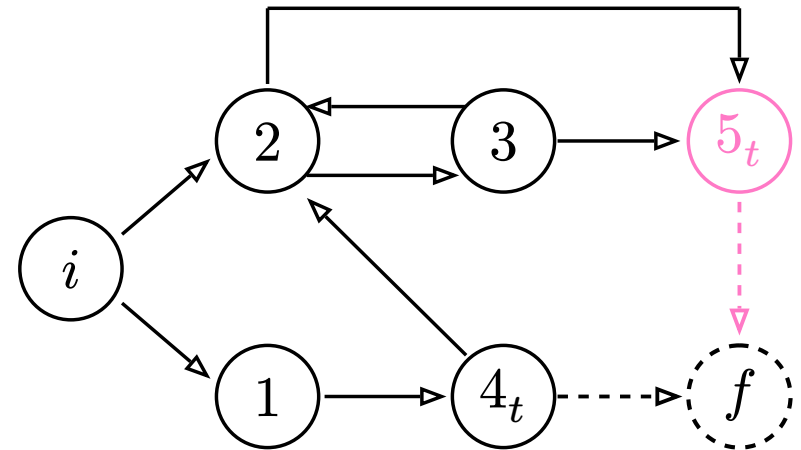
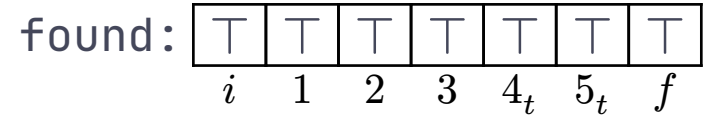
Phase 1 – Call Sequence Generation (Stage 1)



fifo: $[i, 1, 2, 4_t, 3, 5_t]$

cmp : $\{\langle i, 1, 4_t, f \rangle, \langle i, 2, 5_t, f \rangle\}$

inc : $\{\langle i, 1, 4_t, 2 \rangle, \langle i, 2, 3, 2 \rangle, \langle i, 2, 3, 5_t \rangle\}$

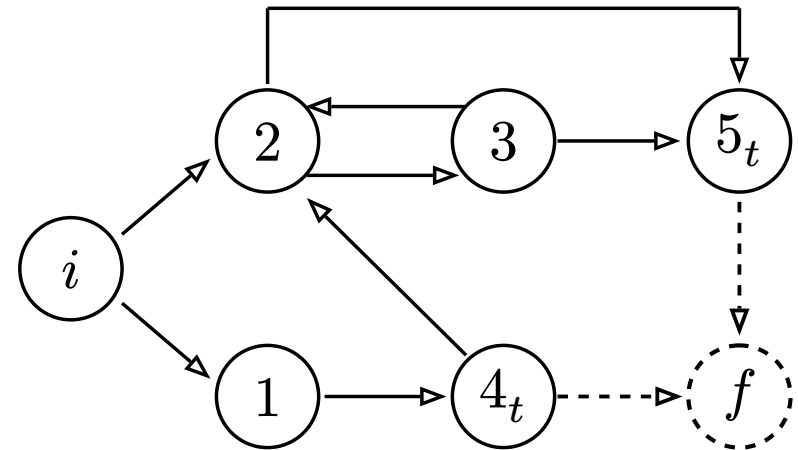
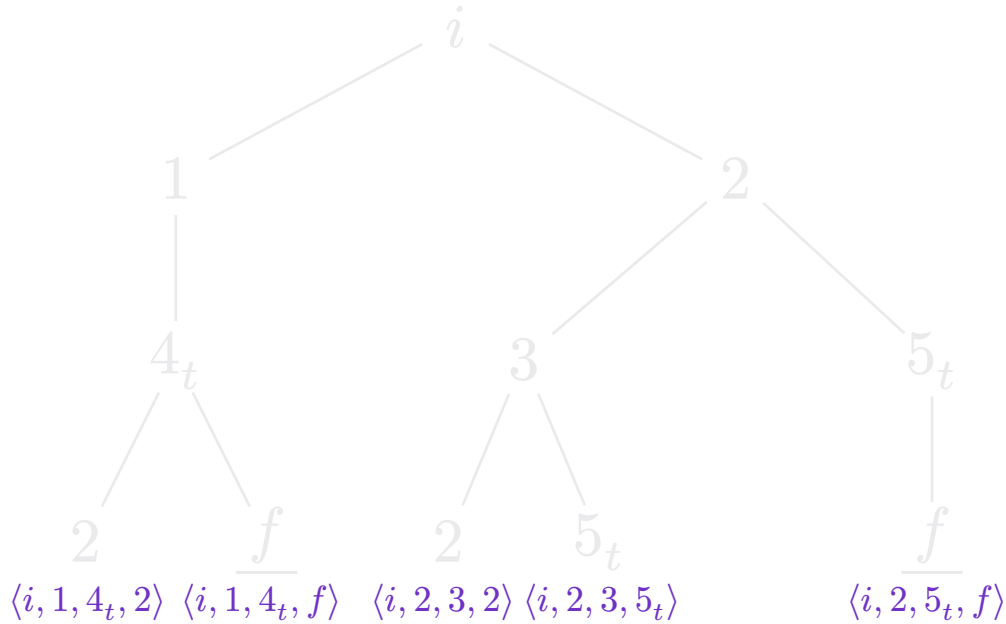


Expanding state 5_t

Final state f has been found

Add path to **cmp**

Phase 1 – Call Sequence Generation (Stage 1)



By the end of stage 1 we know:

1. all states have been explored
2. all transitions belong to some path in **cmp** \cup **inc**
3. we have the shortest paths from *i* to all nodes

fifo: [*i*, 1, 2, 4_t, 3, 5_t]

cmp : { $\langle i, 1, 4_t, f \rangle$, $\langle i, 2, 5_t, f \rangle$ }

inc : { $\langle i, 1, 4_t, 2 \rangle$, $\langle i, 2, 3, 2 \rangle$, $\langle i, 2, 3, 5_t \rangle$ }

Phase 1 – Call Sequence Generation

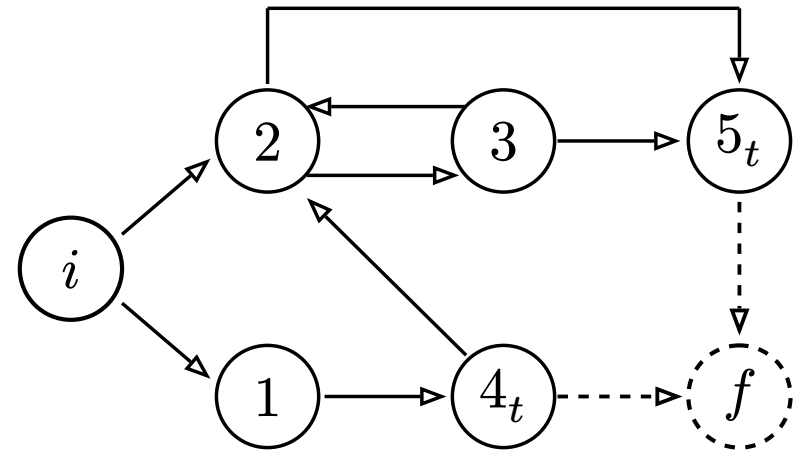
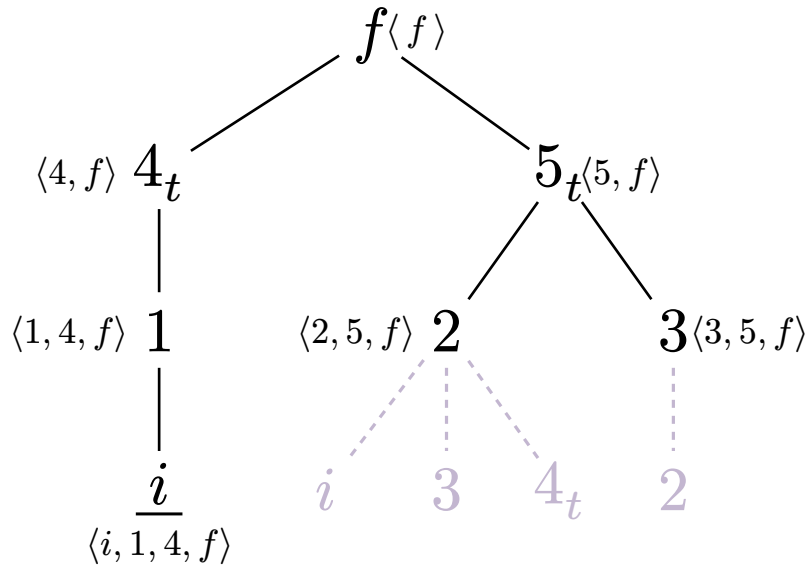
Custom **coverage-guided** algorithm based on breadth-first search (BFS).

The algorithm works in three stages:

1. Collect paths from i , classifying them as complete (reaching f) or incomplete
2. Collect a path from each state to f
3. Complete the incomplete paths



Phase 1 – Call Sequence Generation (Stage 2)



- Working backwards: start from f
- Store a single shortest path from f to all nodes
- By the end of this stage, we know how to complete paths from all nodes to f

Phase 1 – Call Sequence Generation

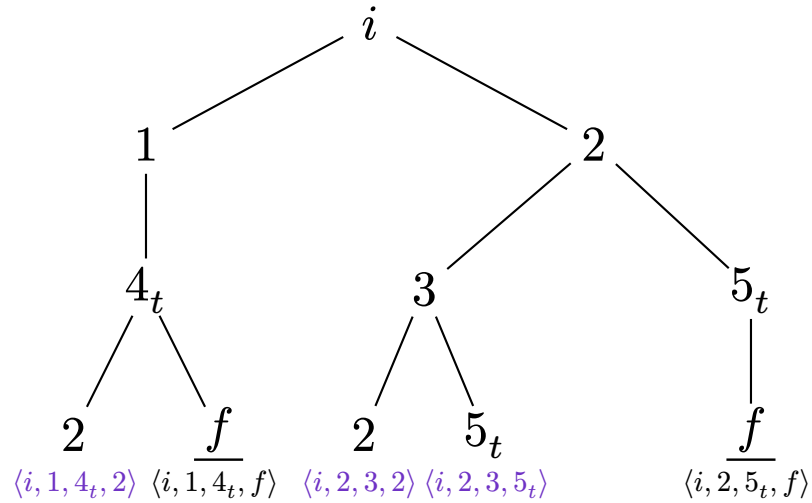
Custom **coverage-guided** algorithm based on breadth-first search (BFS).

The algorithm works in three stages:

1. Collect paths from i , classifying them as complete (reaching f) or incomplete ✓
2. Collect a path from each state to f ✓
3. Complete the incomplete paths

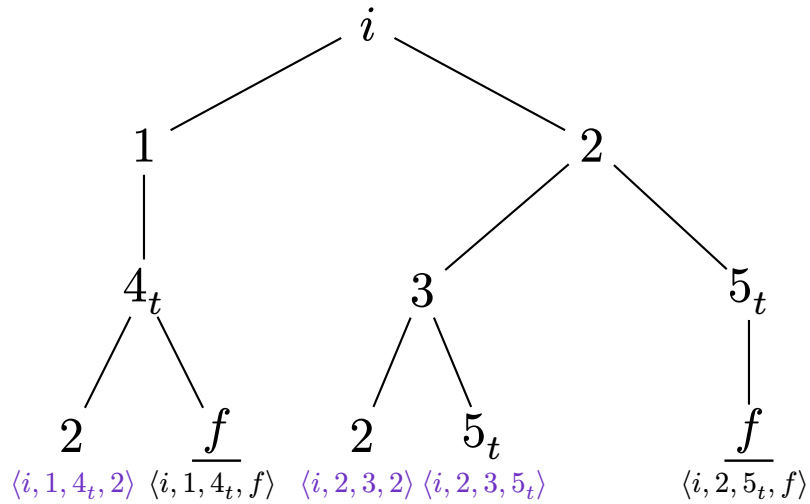
Phase 1 – Call Sequence Generation (Stage 3)

Incomplete paths from Stage 1:

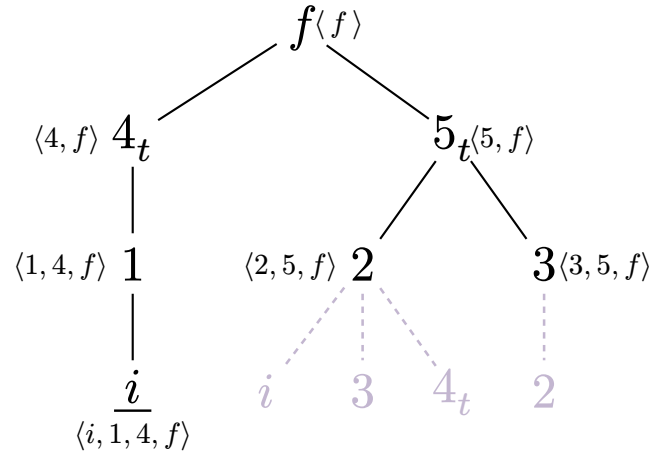


Phase 1 – Call Sequence Generation (Stage 3)

Incomplete paths from Stage 1:

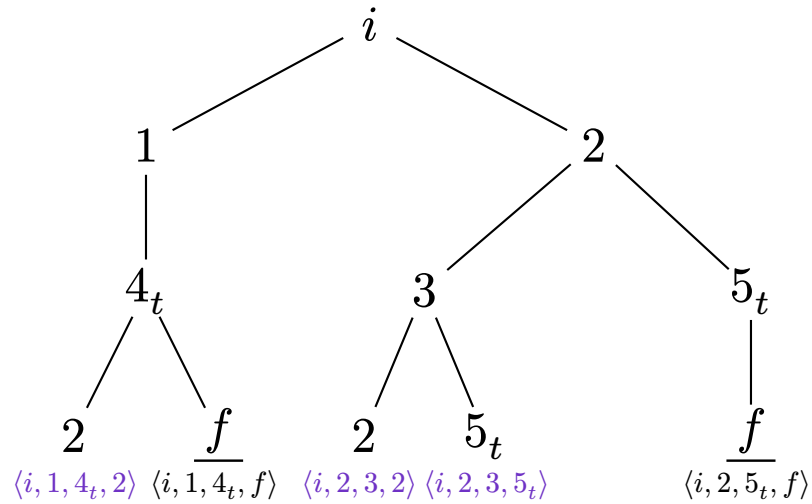


Results from Stage 2:

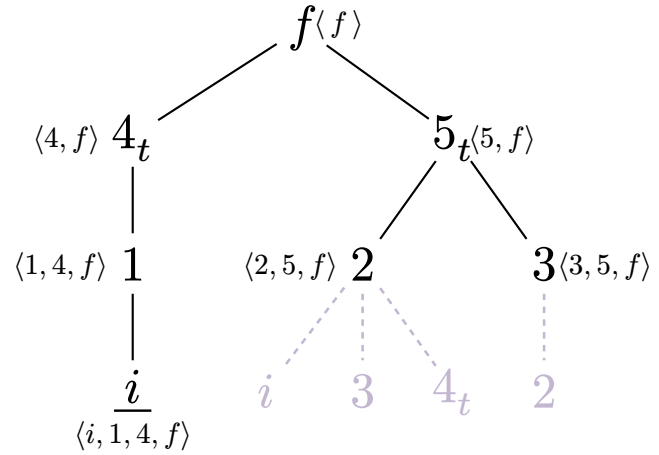


Phase 1 – Call Sequence Generation (Stage)

Incomplete paths from Stage 1:



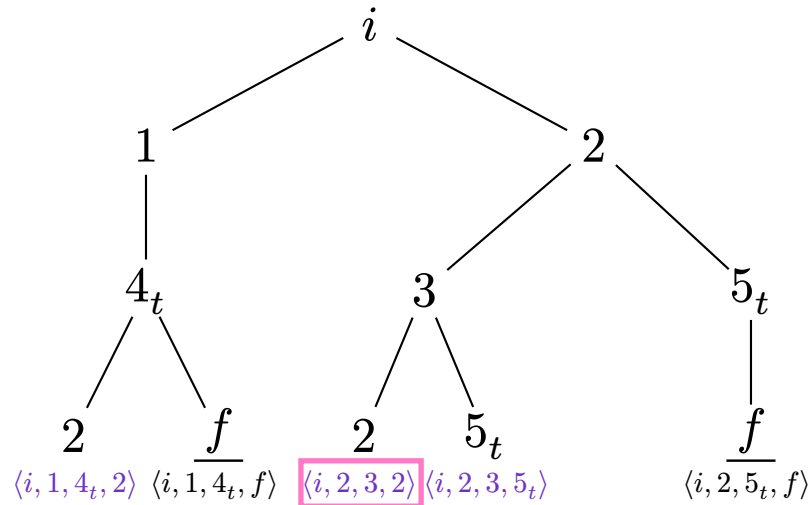
Result from Stage 2:



Finish the incomplete paths from stage 1 with the paths from stage 2.

Phase 1 – Call Sequence Generation (Stage 3)

Incomplete paths from Stage 1:



Result from Stage 2:

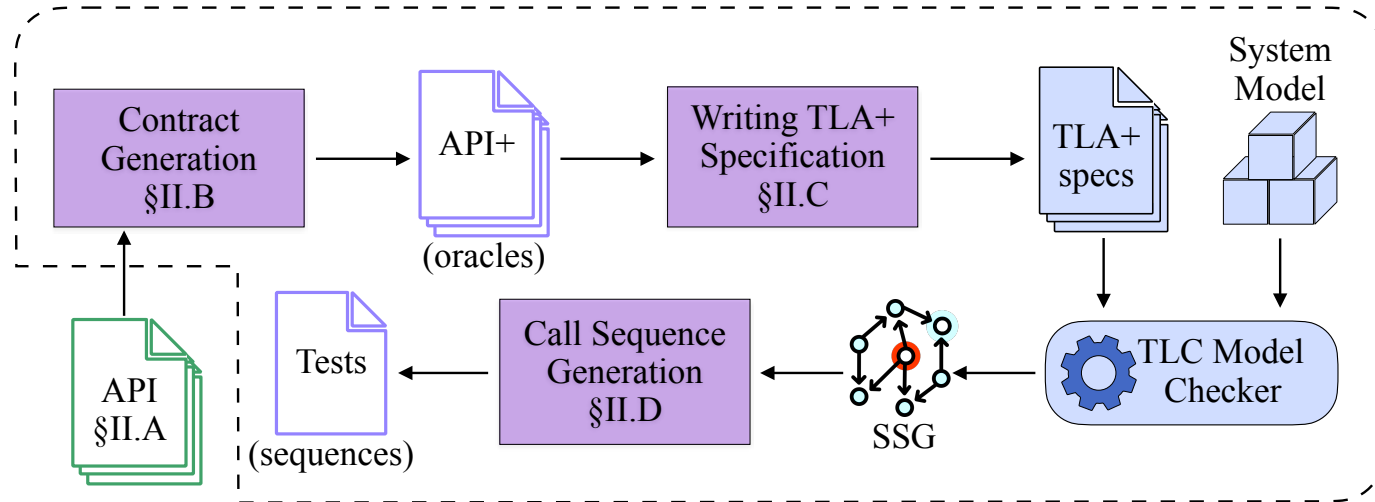


Finish incomplete paths from stage 1 with the paths from stage 2.

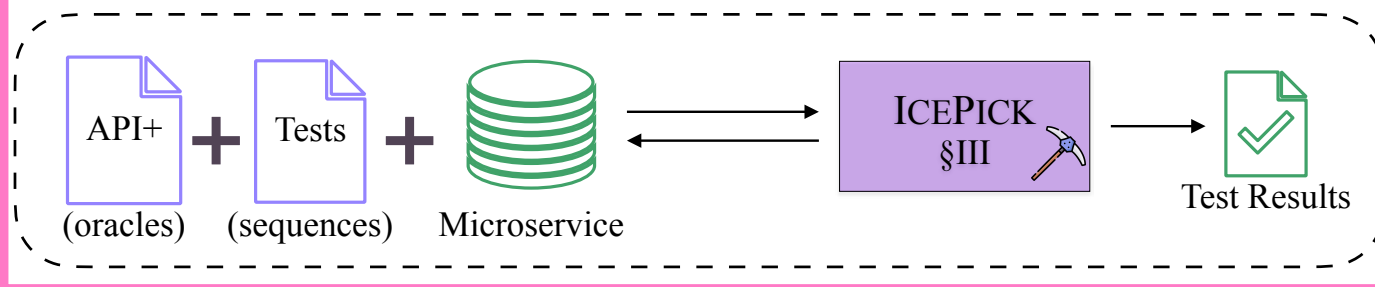
$$\langle i, 2, 3, 2 \rangle \oplus \langle 2, 5, f \rangle = \langle i, 2, 3, 2, 5, f \rangle$$

Phase 2 – Testing

ICEPICK Phase 1 - Specification pre-processing

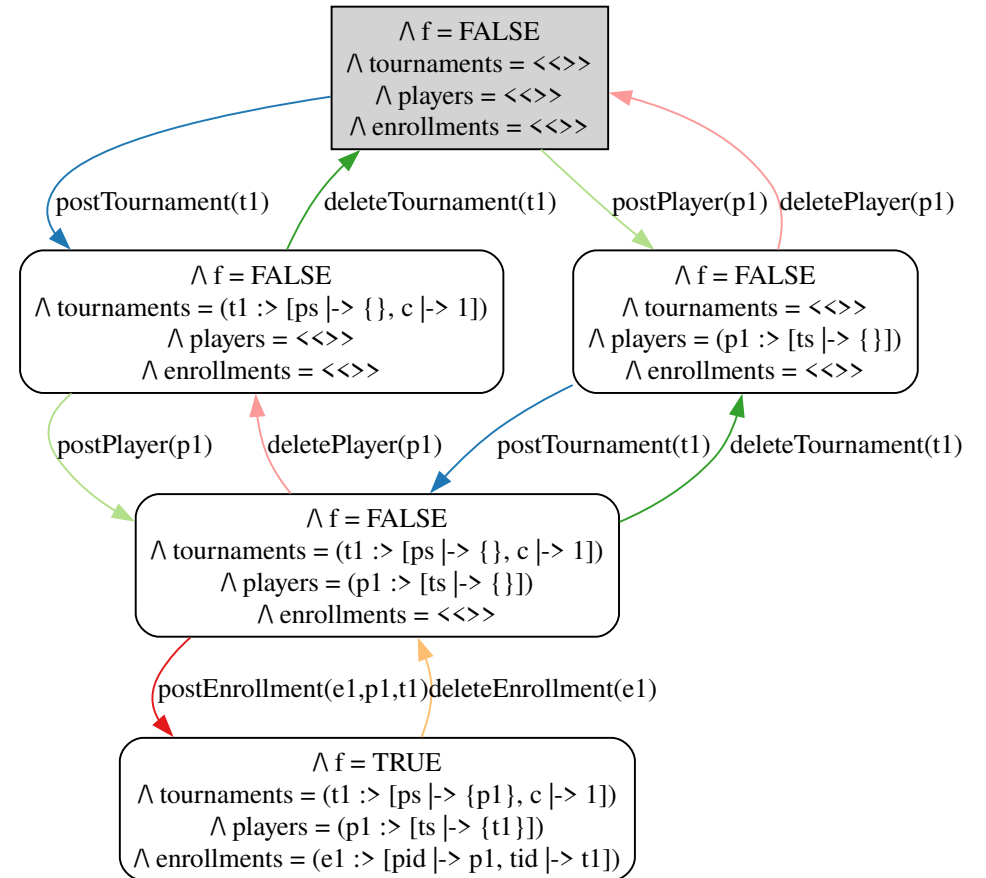


ICEPICK Phase 2 - Testing



Phase 2 – Testing

SSG resulting from running TLC with the Tournaments specification and a model with one model-value per resource type.

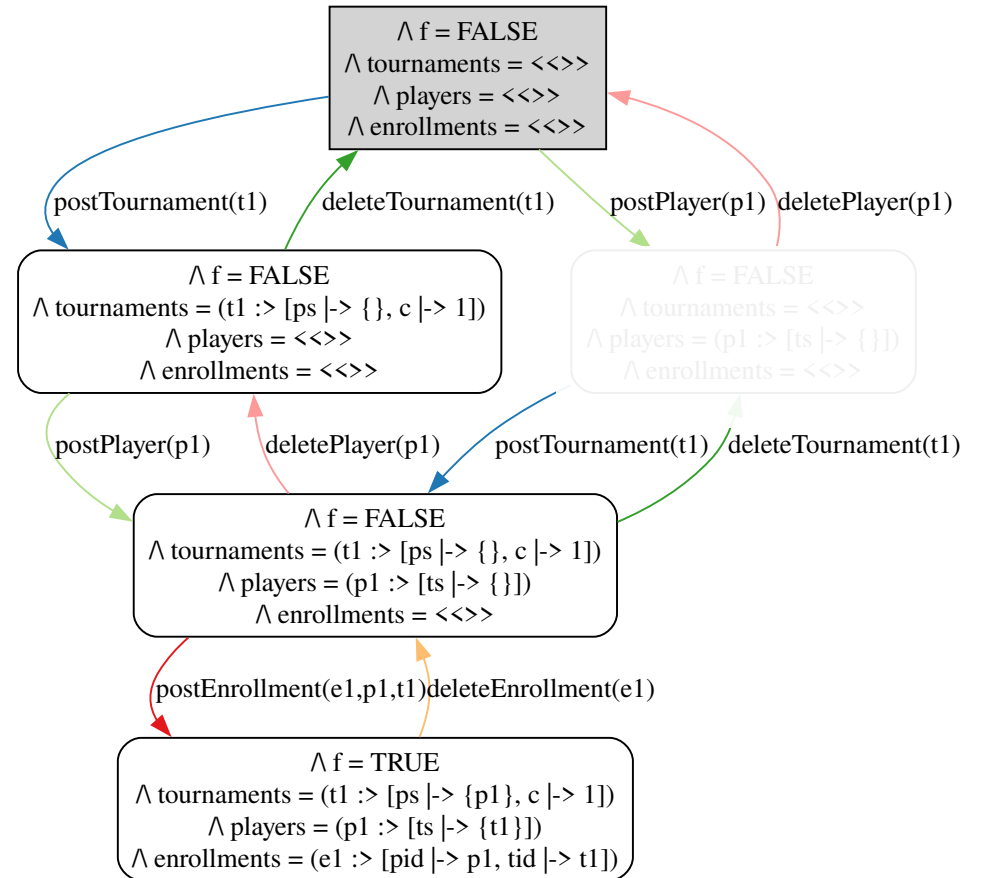


Phase 2 – Testing

SSG resulting from running TLC with the Tournaments specification and a model with one model-value per resource type.

Let's focus on a single path:

**postTournament(t1) → postPlayer(p1) →
postEnrolment(e1) → deleteEnrolment(e1)**



Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

```
data ← verb = post ?  
  Generate(op) : Recycle(op)  
  
invs ← eval(op.invs)  
precond ← eval(req, data)  
  
res ← verb = post ?  
  post(op, data) : delete(op, data)  
  
if res.code ≠ 5xx  
  postcond ← eval(ens, data, res)  
  
invs ← eval(op.invs)  
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postTournament

- Generate → t1

Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

```
data ← verb = post ?
  Generate(op) : Recycle(op)
invs ← eval(op.invs)
precond ← eval(req, data)

res ← verb = post ?
  post(op, data) : delete(op, data)
if res.code ≠ 5xx
  postcond ← eval(ens, data, res)
invs ← eval(op.invs)
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postTournament

- Generate → t1
- Evaluate invariants and preconds with t1

Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

```
data ← verb = post ?
  Generate(op) : Recycle(op)
invs ← eval(op.invs)
precond ← eval(req, data)
res ← verb = post ?
  post(op, data) : delete(op, data)
if res.code ≠ 5xx
  postcond ← eval(ens, data, res)
invs ← eval(op.invs)
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postTournament

- Generate → t1
- Evaluate invariants and preconds with t1
- Make the HTTP request with t1

Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

```
data ← verb = post ?
  Generate(op) : Recycle(op)
invs ← eval(op.invs)
precond ← eval(req, data)

res ← verb = post ?
  post(op, data) : delete(op, data)
if res.code ≠ 5xx
  postcond ← eval(ens, data, res)
invs ← eval(op.invs)

result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postTournament

- Generate → t1
- Evaluate invariants and preconds with t1
- Make the HTTP request with t1
- Evaluate postconds (when the server does not crash) and invariants

Phase 2 – Testing Algorithm

`postTournament(t1)` → `postPlayer(p1)` → `postEnrolment(e1)` → `deleteEnrolment(e1)`

```
data ← verb = post ?  
  Generate(op) : Recycle(op)  
invs ← eval(op.invs)  
precond ← eval(req, data)  
res ← verb = post ?  
  post(op, data) : delete(op, data)  
if res.code ≠ 5xx  
  postcond ← eval(ens, data, res)  
invs ← eval(op.invs)  
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postTournament

- Generate → t1
- Evaluate invariants and preconds with t1
- Make the HTTP request with t1
- Evaluate postconds (when the server does not crash) and invariants
- Oracle: based on invariants, preconds, postconds and response codes

Phase 2 – Testing Algorithm

postTournament(t1) → postPlayer(p1) → postEnrolment(e1) → deleteEnrolment(e1)

```
data ← verb = post ?  
  Generate(op) : Recycle(op)  
invs ← eval(op.invs)  
precond ← eval(req, data)  
  
res ← verb = post ?  
  post(op, data) : delete(op, data)  
  
if res.code ≠ 5xx  
  postcond ← eval(ens, data, res)  
  
invs ← eval(op.invs)  
  
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postPlayer

- Exactly the same procedure, since it's a post operation.

Phase 2 – Testing Algorithm

postTournament(t1) → postPlayer(p1) → **postEnrolment(e1)** → deleteEnrolment(e1)

```
data ← verb = post ?  
  Generate(op) : Recycle(op)  
invs ← eval(op.invs)  
precond ← eval(req, data)  
  
res ← verb = post ?  
  post(op, data) : delete(op, data)  
  
if res.code ≠ 5xx  
  postcond ← eval(ens, data, res)  
  
invs ← eval(op.invs)  
  
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: postEnrolment

- Exactly the same procedure, since it's a post operation.

Phase 2 – Testing Algorithm

postTournament(t1) → postPlayer(p1) → postEnrolment(e1) → deleteEnrolment(e1)

```
data ← verb = post ?  
  Generate(op) : Recycle(op)  
  
invs ← eval(op.invs)  
precond ← eval(req, data)  
  
res ← verb = post ?  
  post(op, data) : delete(op, data)  
  
if res.code ≠ 5xx  
  postcond ← eval(ens, data, res)  
  
invs ← eval(op.invs)  
result ← invs ∧ precond ∧ postcond ∧ res.code
```

Testing: deleteEnrolment

- Instead of generating data, we're recycling.

Evaluation Results

We evaluated IcePick on:

- EvoMaster Benchmark (Features Service, E-Commerce API)
- Petstore (canonical example for OAS)
- Tournaments Management System (both with and without fault-injection)

SYSTEMS UNDER TEST DESCRIPTION.

System	Endpoints	Operations	Parameters	Schemas
Tournaments	11	19	15	4
Swagger-Petstore	13	19	17	6
E-Commerce	22	28	33	25
Features-Service	10	16	33	8

Evaluation Results

Some errors were artificially injected, others we found organically.

TEST RESULTS.

	Model	Exec. Time	WARN	ERR	NOT_TESTED
TOUR.	p1_t1_e1	00min 02s	0	0	0
	p2_t2_e2	02min 20s	0	0	0
	p1_t1_e1_ext	00min 14s	0	0	0
T. ERR	p1_t1_e1_ERR	00min 02s	8	5	0
	p2_t2_e2_ERR	02min 44s	1658	459	49
	p1_t1_e1_ERR_ext	00min 17s	7	5	0
PETS.	u1_p1_o1	00min 01s	0	0	3
	u2_p2_o2	00min 11s	0	0	72
FS.	p1_c1_f1	< 01s	0	97	50

Evaluation Results

Some errors were artificially injected, others were found organically.

IcePick successfully found all injected errors and REST-compliance issues with some benchmark systems.

TEST RESULTS.

	Model	Exec. Time	WARN	ERR	NOT_TESTED
TOUR.	p1_t1_e1	00min 02s	0	0	0
	p2_t2_e2	02min 20s	0	0	0
	p1_t1_e1_ext	00min 14s	0	0	0
T. ERR	p1_t1_e1_ERR	00min 02s	8	5	0
	p2_t2_e2_ERR	02min 44s	1658	459	49
	p1_t1_e1_ERR_ext	00min 17s	7	5	0
PETS.	u1_p1_o1	00min 01s	0	0	3
	u2_p2_o2	00min 11s	0	0	72
FS.	p1_c1_f1	< 01s	0	97	50

Evaluation Results

Some errors were artificially injected, others were found organically.

IcePick successfully found all injected errors and REST-compliance issues with some benchmark systems.

Detected Referential Integrity violation:

1. postTournament(t1)
2. postPlayer(p1)
3. postEnrolment(e1)
4. deleteEnrolment(e1)

TEST RESULTS.

	Model	Exec. Time	WARN	ERR	NOT_TESTED
TOUR.	p1_t1_e1	00min 02s	0	0	0
	p2_t2_e2	02min 20s	0	0	0
	p1_t1_e1_ext	00min 14s	0	0	0
T. ERR	p1_t1_e1_ERR	00min 02s	8	5	0
	p2_t2_e2_ERR	02min 44s	1658	459	49
	p1_t1_e1_ERR_ext	00min 17s	7	5	0
PETS.	u1_p1_o1	00min 01s	0	0	3
	u2_p2_o2	00min 11s	0	0	72
FS.	p1_c1_f1	< 01s	0	97	50

Future Work

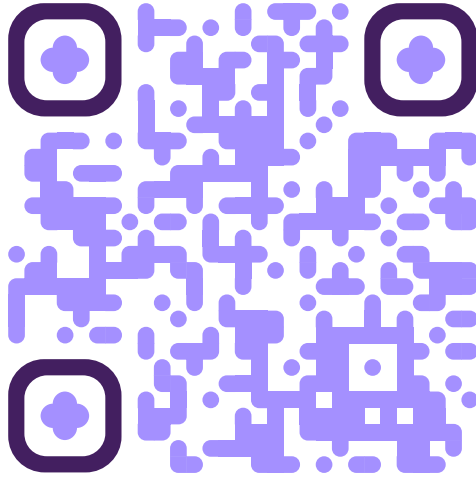
- Automate the TLA+ specification generation
 - Glacier → TLA+
- Deal with partially compliant OAS specifications
- Test on real-world systems

TEST RESULTS.

	Model	Exec. Time	WARN	ERR	NOT_TESTED
TOUR.	p1_t1_e1	00min 02s	0	0	0
	p2_t2_e2	02min 20s	0	0	0
	p1_t1_e1_ext	00min 14s	0	0	0
T. ERR	p1_t1_e1_ERR	00min 02s	8	5	0
	p2_t2_e2_ERR	02min 44s	1658	459	49
	p1_t1_e1_ERR_ext	00min 17s	7	5	0
PETS.	u1_p1_o1	00min 01s	0	0	3
	u2_p2_o2	00min 11s	0	0	72
FS.	p1_c1_f1	< 01s	0	97	50

Want to know more?

Our work was accepted at the 2026 edition of the
IEEE International Conference on Software Testing, Verification and Validation (ICST)



Artifact Available & Reviewed

Thank You



Questions?

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github.com/acm-ribeiro